



Finding The Most Preferred Path

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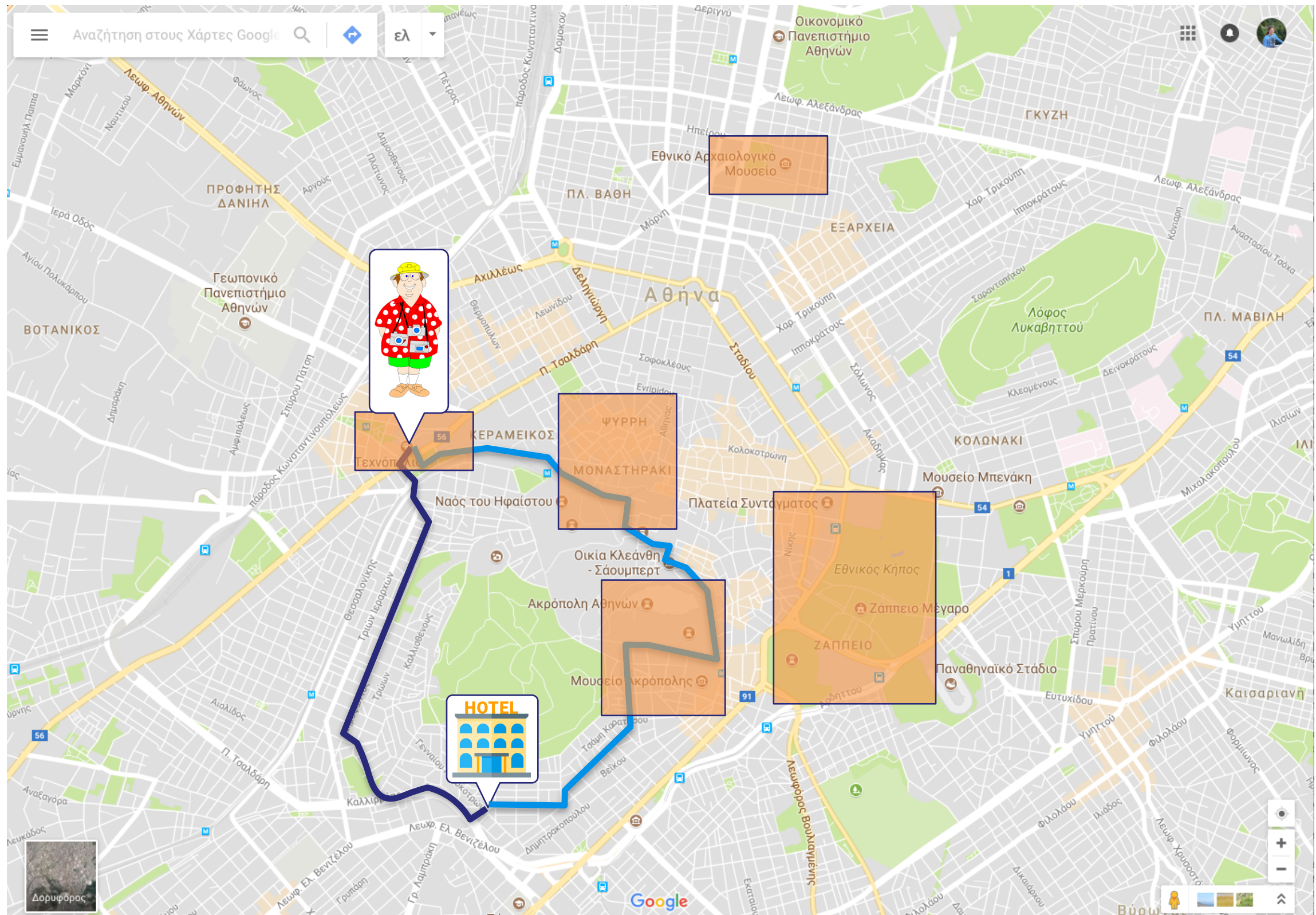
Overview

- Motivating Scenarios
- Preferred Network
- Most Preferred Unrestricted Path
- Most Preferred Near Shortest Path
- Conclusions & Future Work

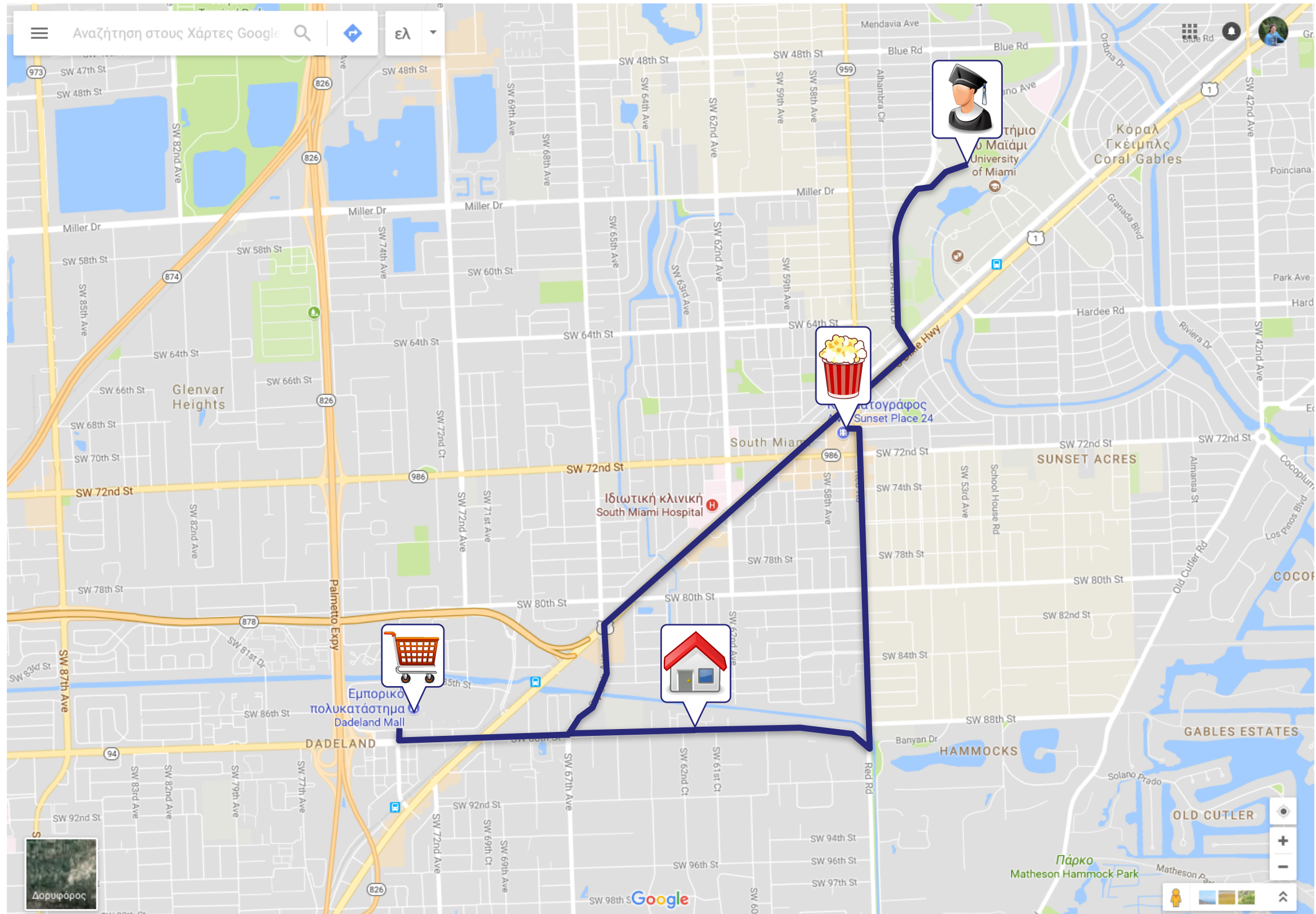
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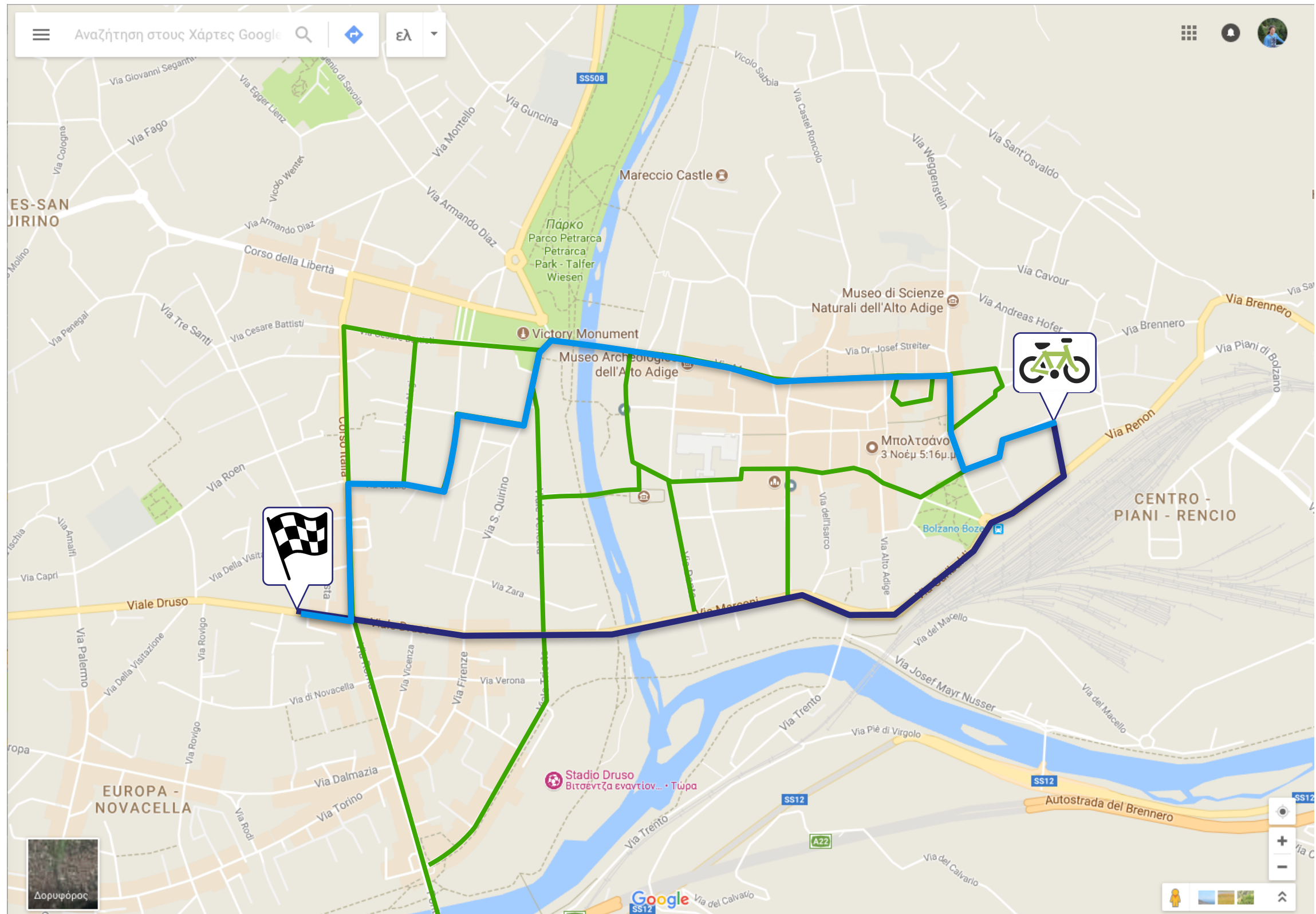
Example 1 - Scenic Route Planning



Example 2 - Familiar Roads



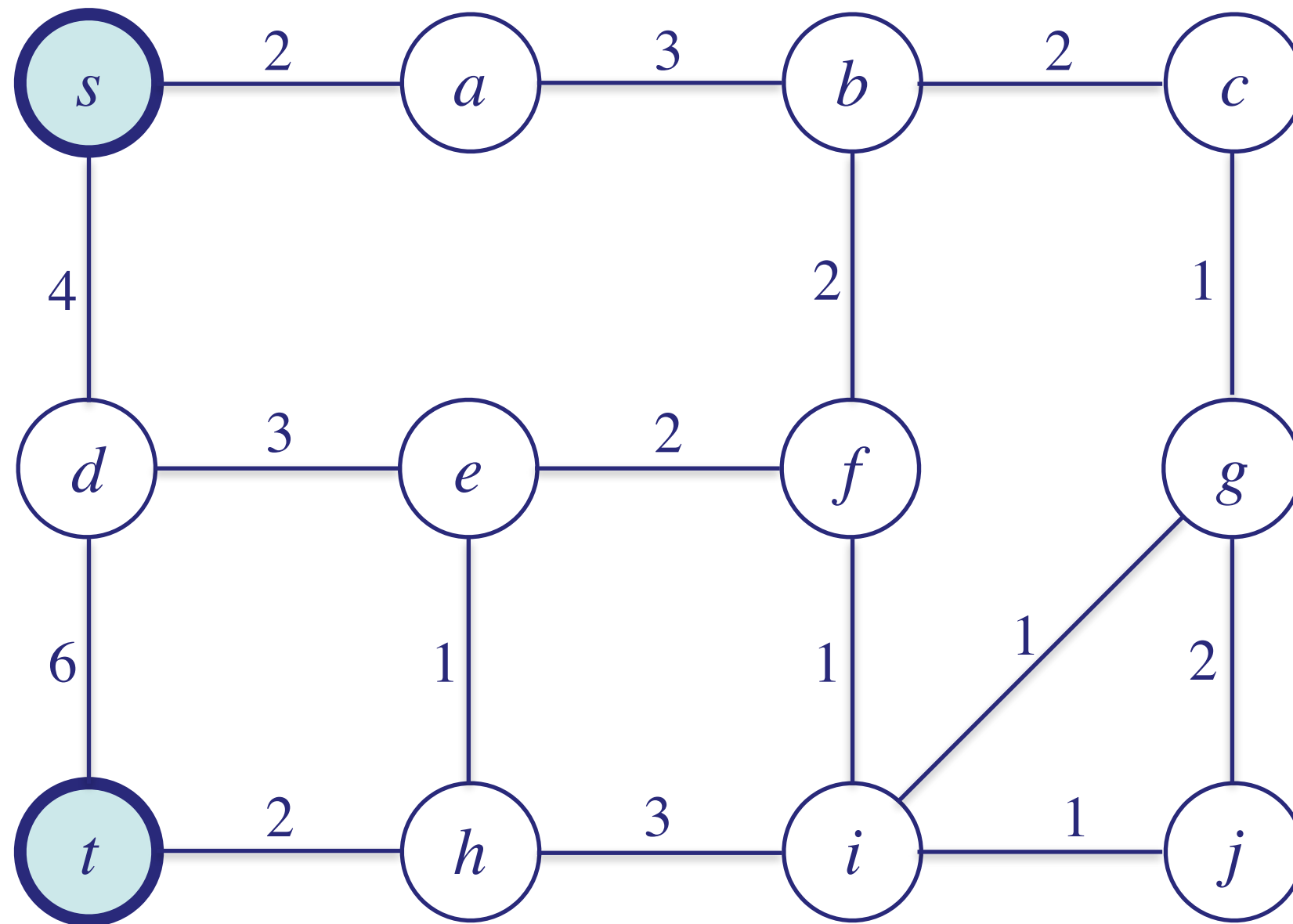
Example 3 - Bicycle Routes



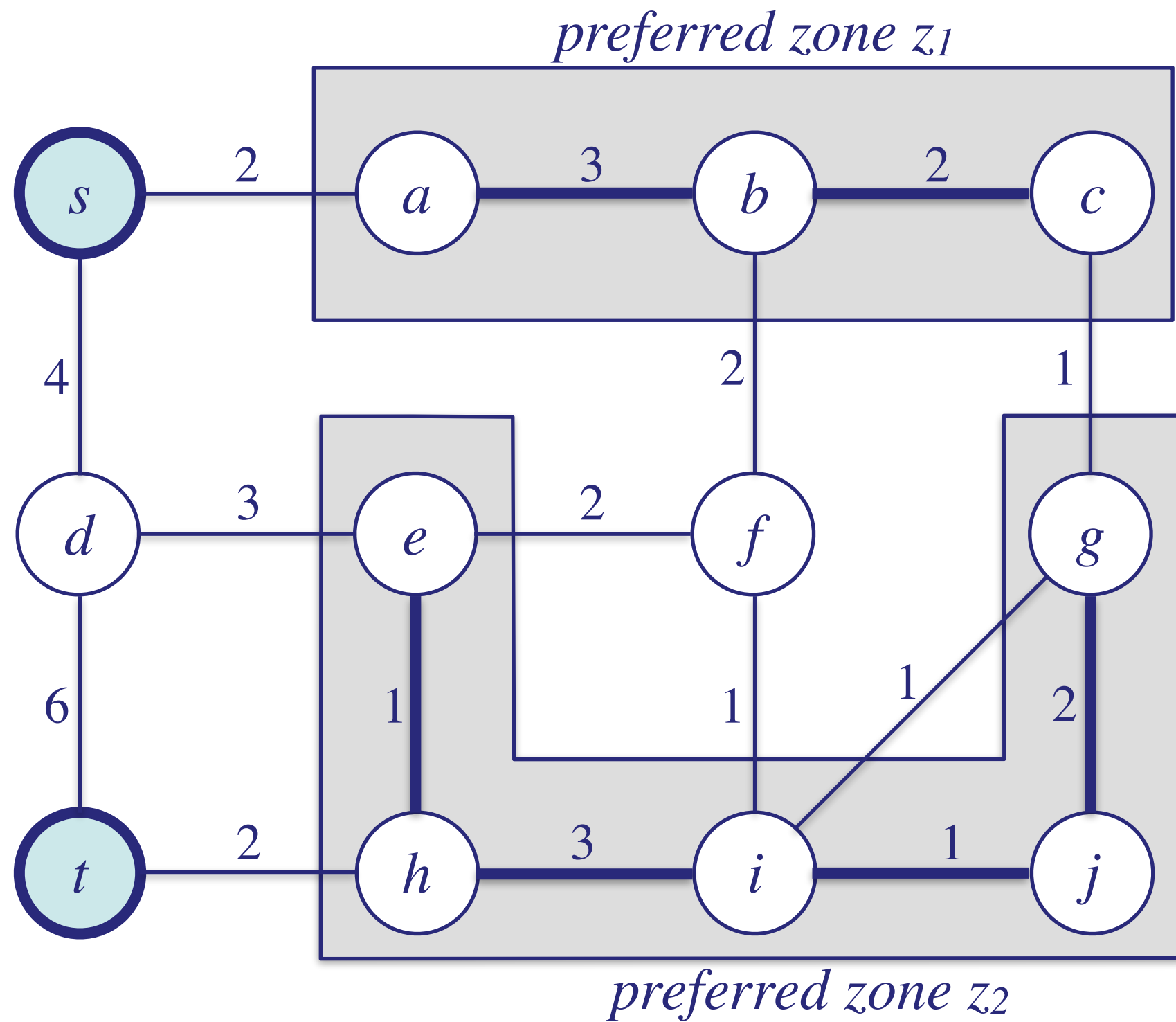
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Road Network



Preferred Network

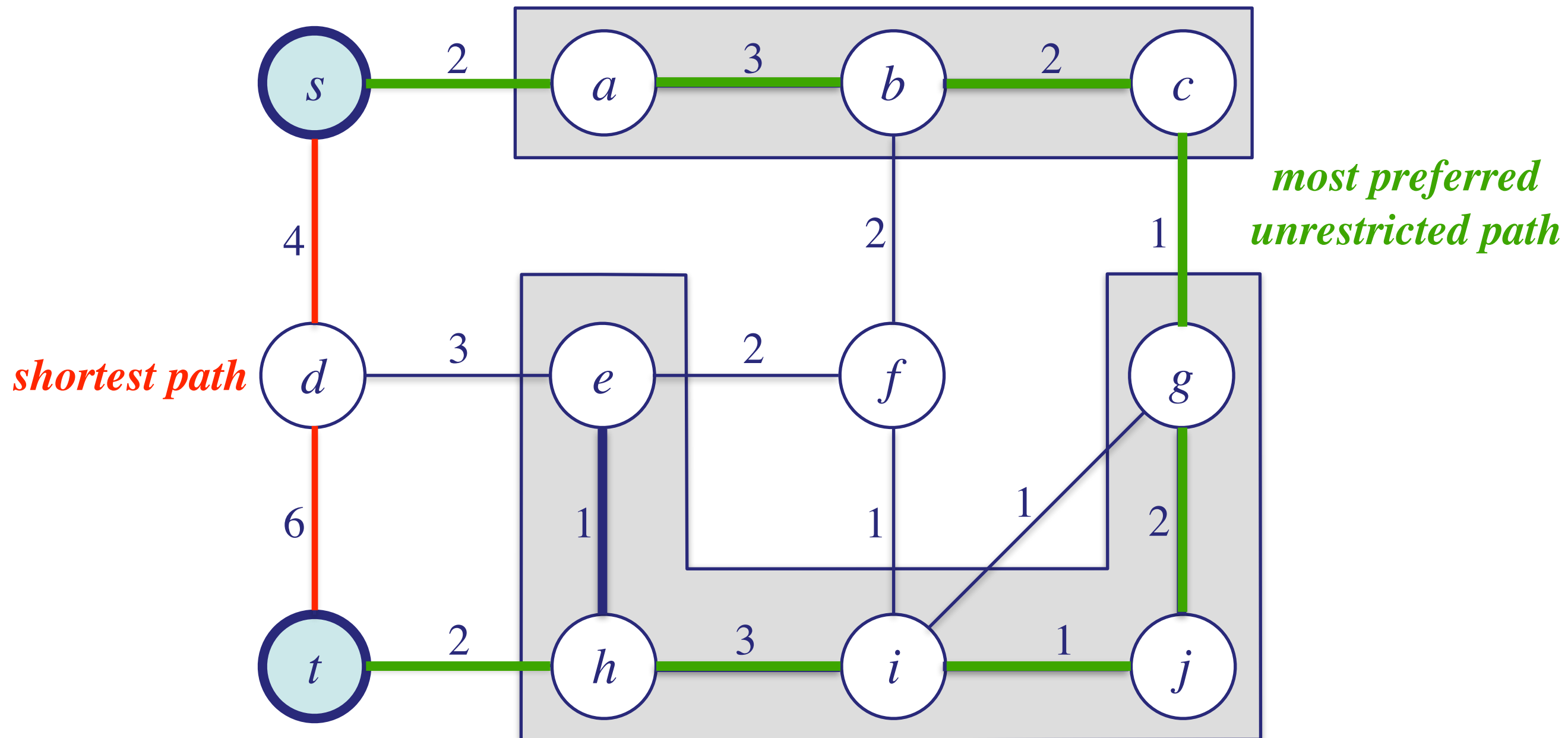


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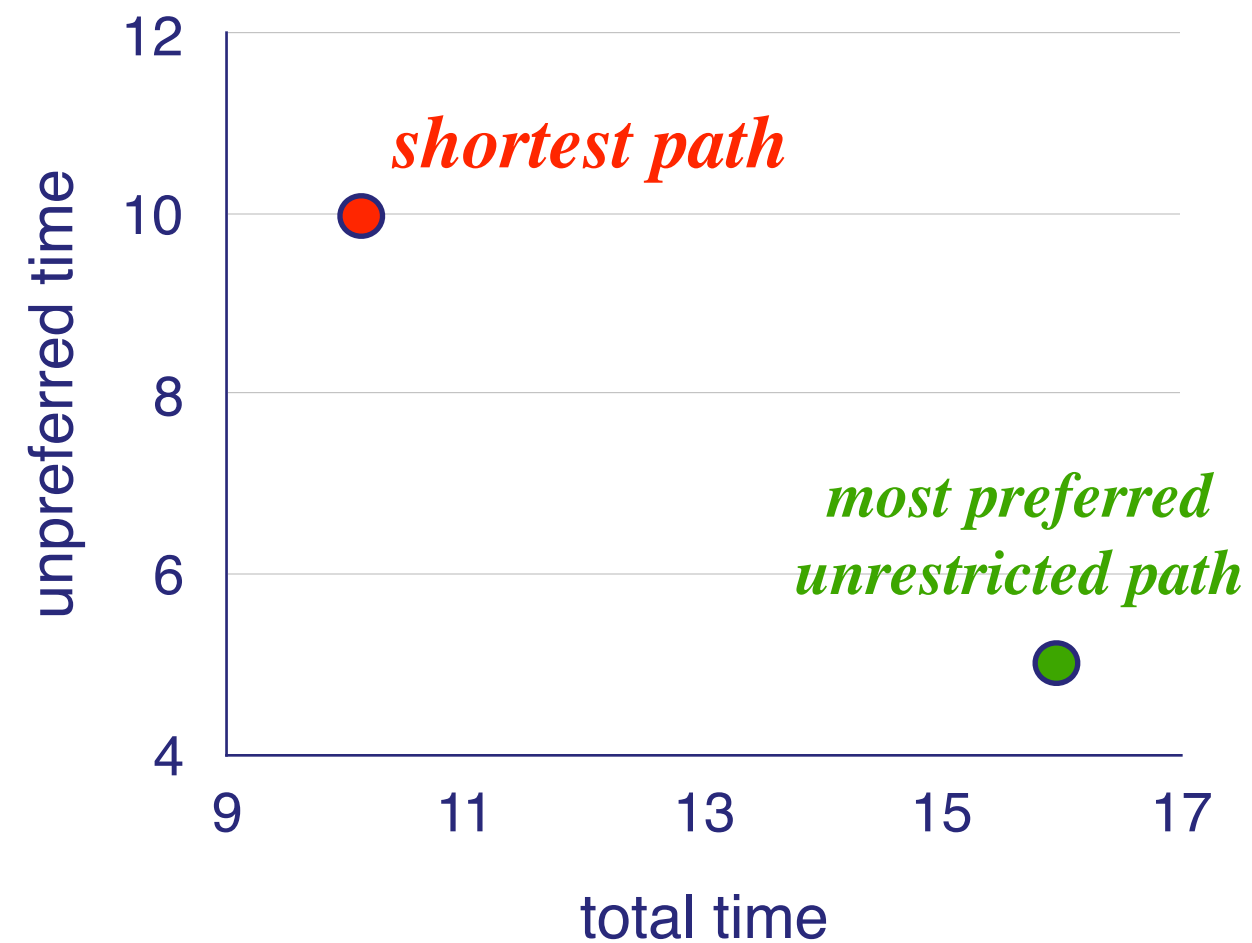
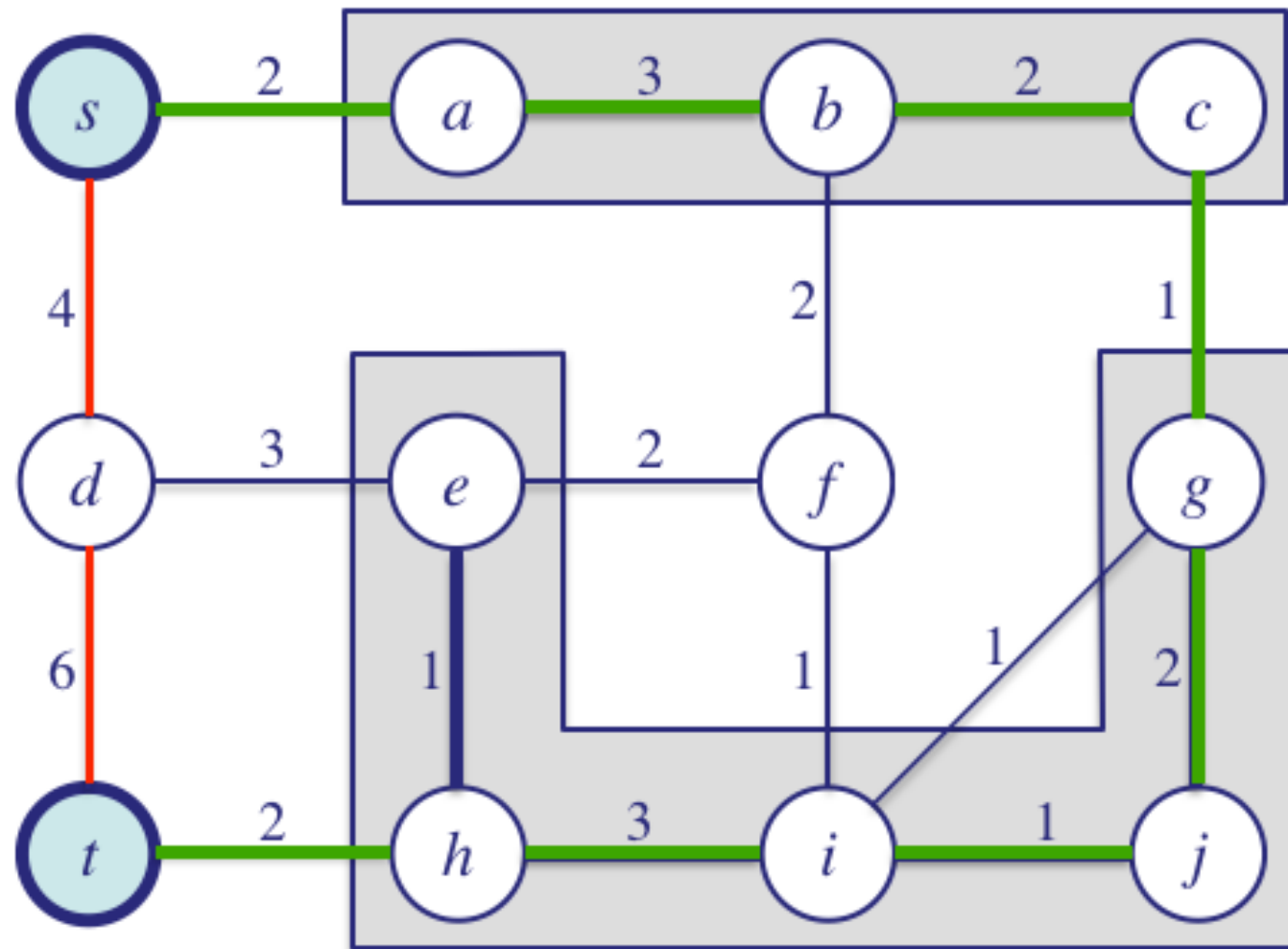
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Most Preferred Unrestricted Path

- Find the path which minimizes the time spent outside the Preferred Network



Most Preferred Unrestricted Path

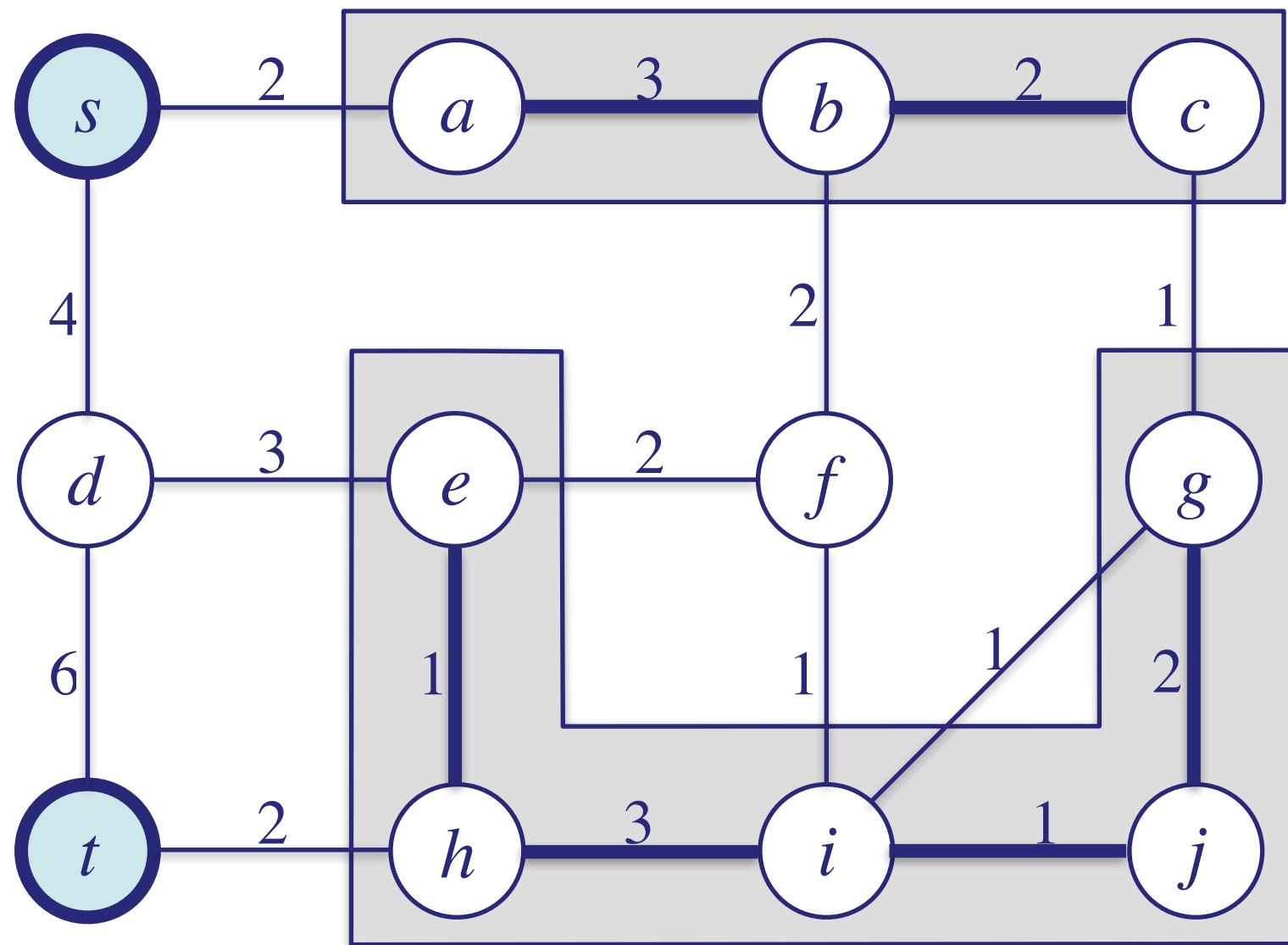


HyperEdges

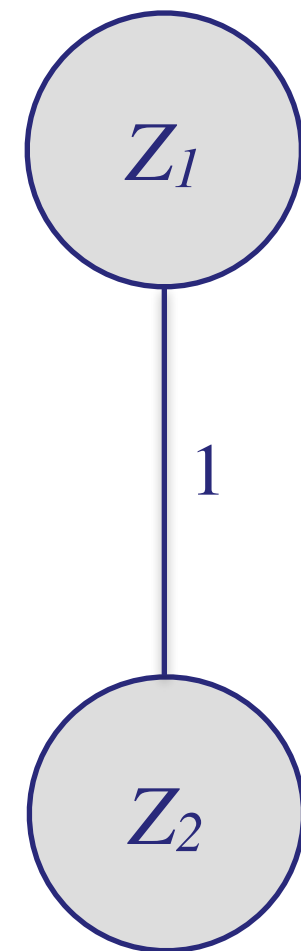
- Introduced by *Aljubayrin et al* in *ICDE 2015* (*Safest Path via Safe Zones*)
- Offline phase: Hypergraph construction
 1. Zones become HyperNodes
 2. HyperEdges are added between HyperNodes/Zones
 3. Weights are determined by shortest paths connecting different zones

HyperEdges

Original Network



Hypergraph

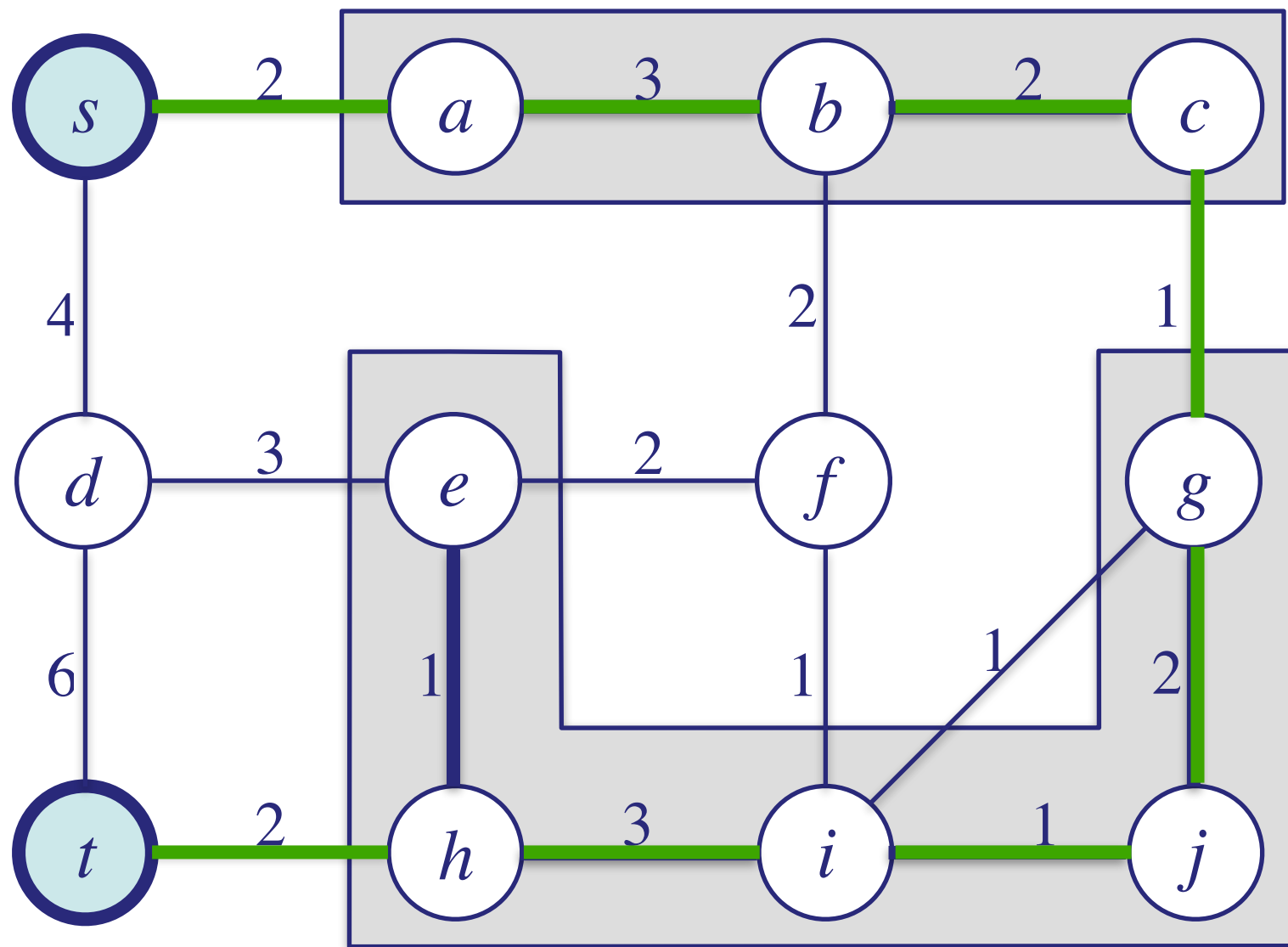


HyperEdges

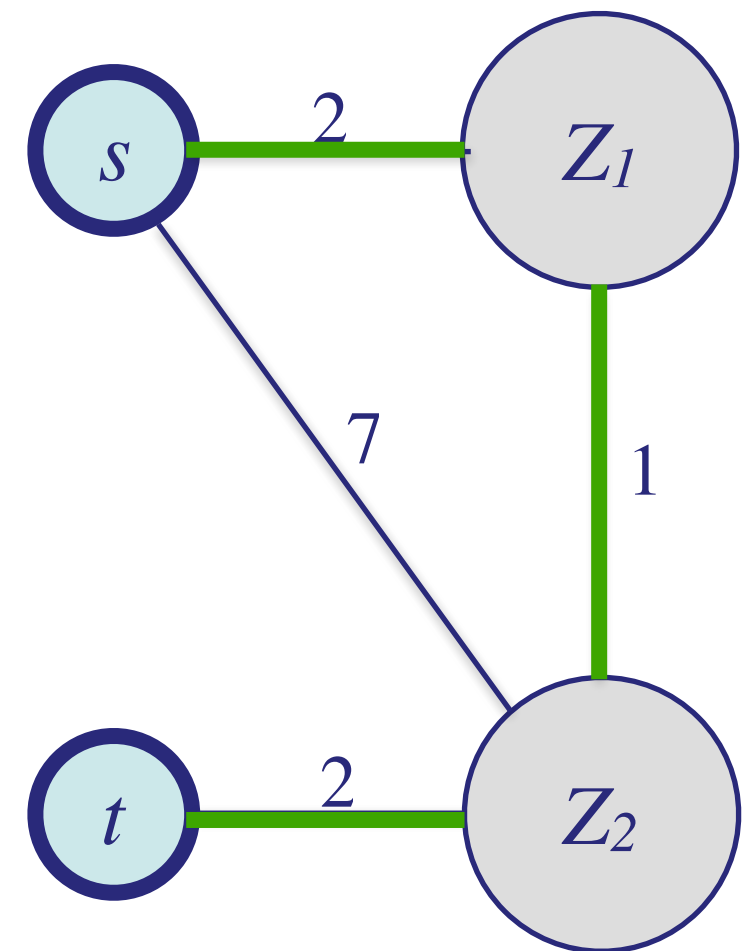
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- Offline phase: Hypergraph construction
 1. Zones become HyperNodes
 2. HyperEdges are added between HyperNodes/Zones
 3. Weights are determined by shortest paths connecting different zones
- Online phase: Query Processing
 1. Add source and target nodes to Hypergraph
 2. Run shortest path query over the Hypergraph

HyperEdges

Original Network



Hypergraph

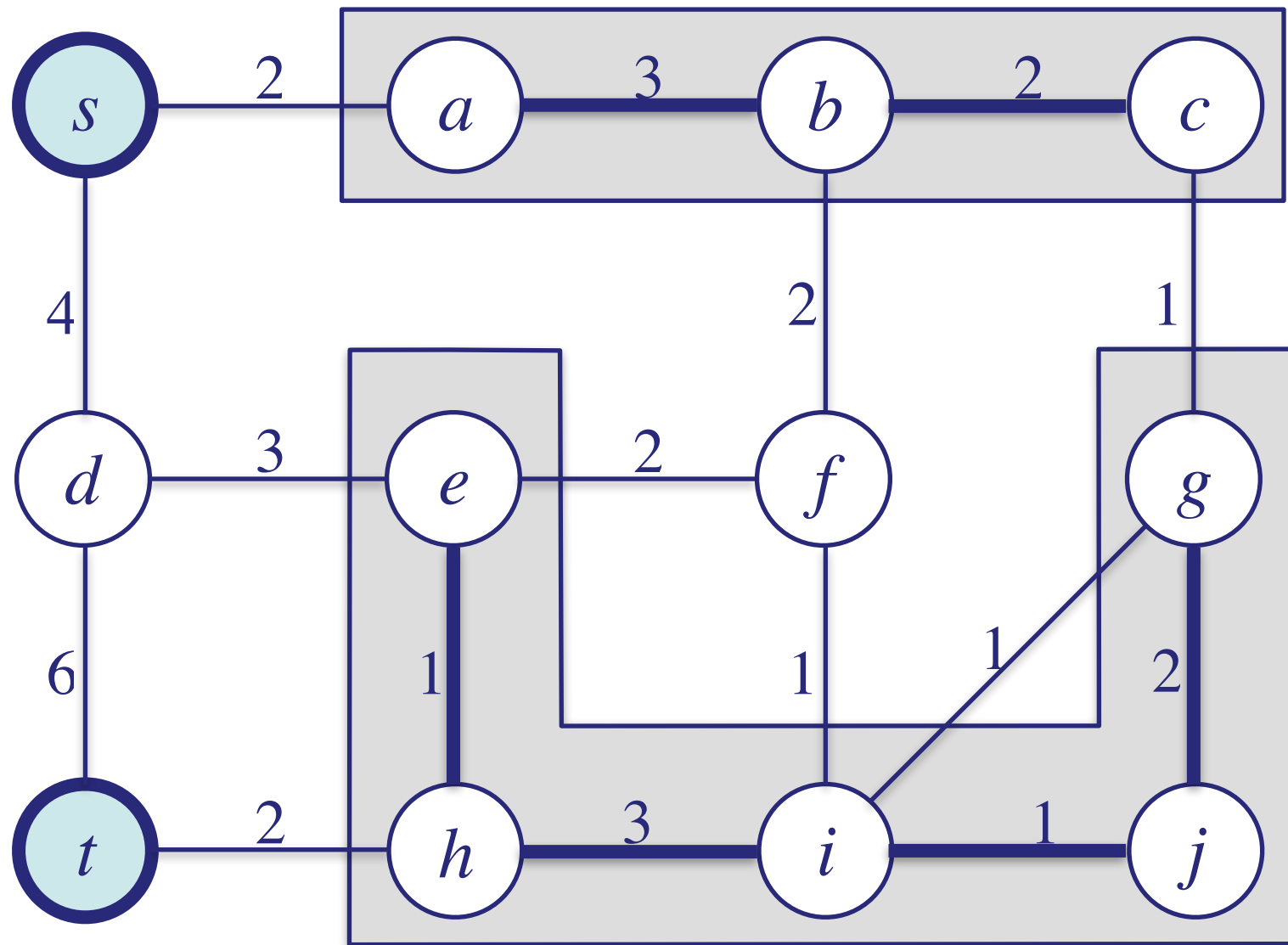


Compressed Network

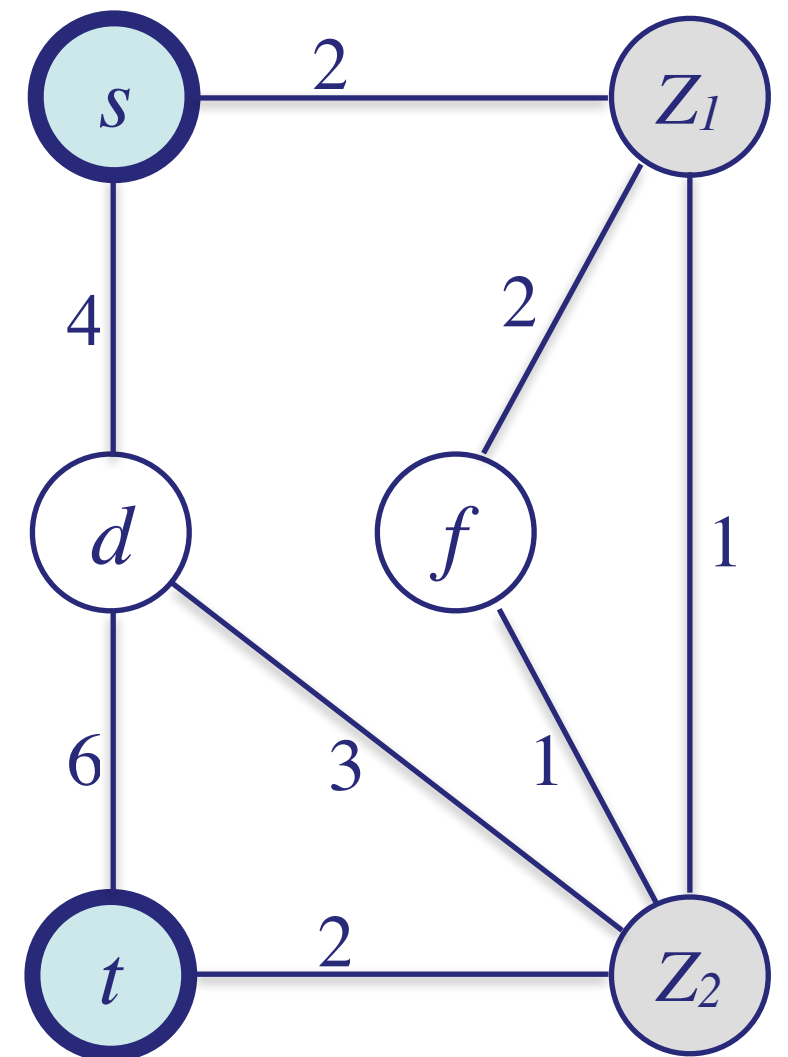
- Offline phase: Build Compressed Network
 1. Replace every zone with a node
 2. Add edges between each new node and each node previously connected with the associated zone

Compressed Network

Original Network



Compressed Network



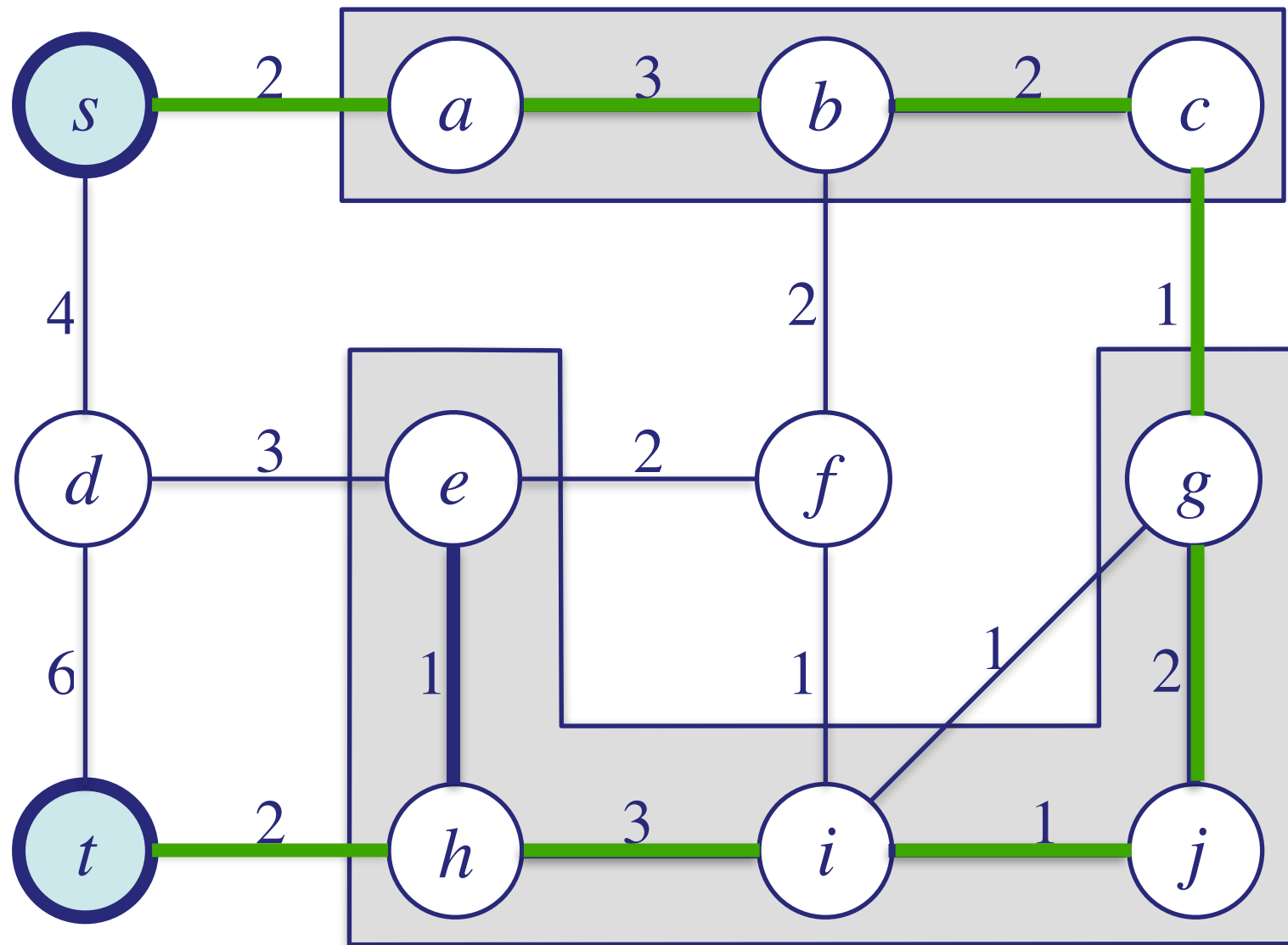
Compressed Network

- Offline phase: Build Compressed Network
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- Online phase: Query Processing

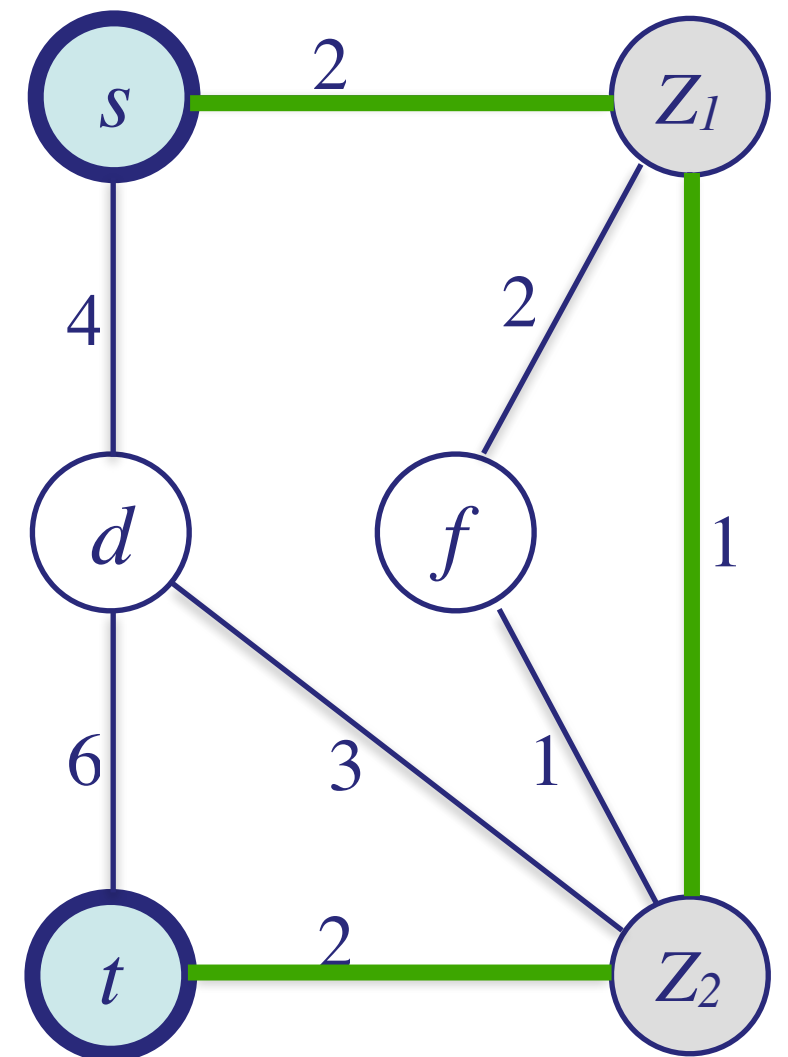
Run shortest path query over the Compressed Network

Compressed Network

Original Network



Compressed Network

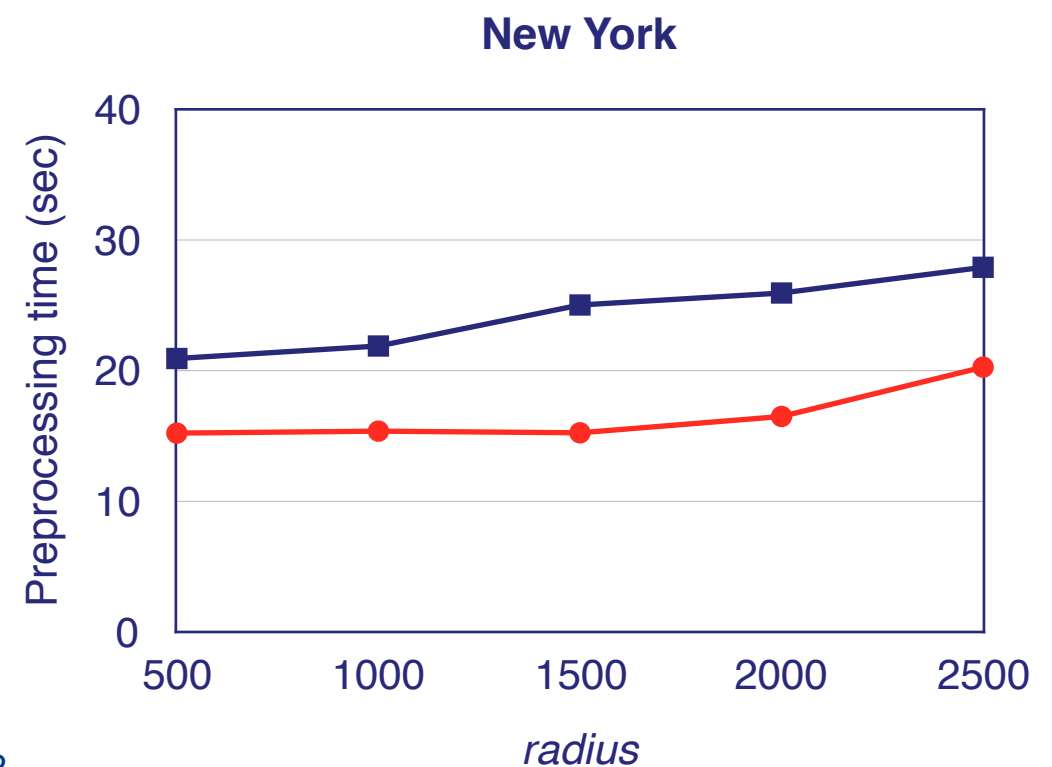
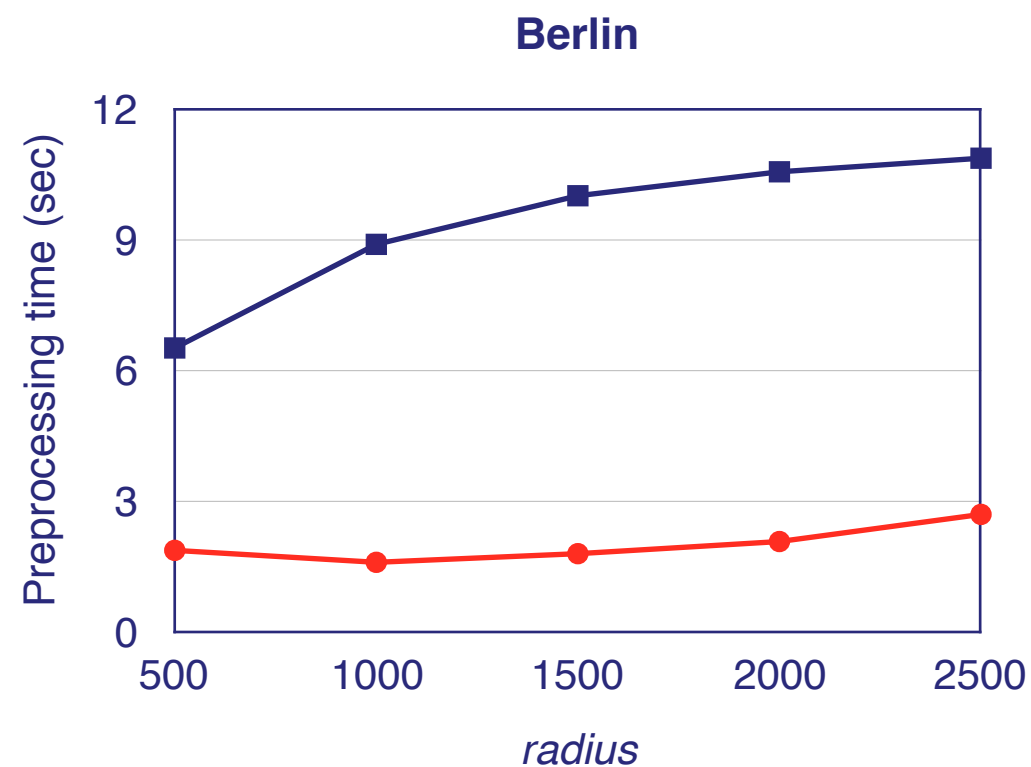
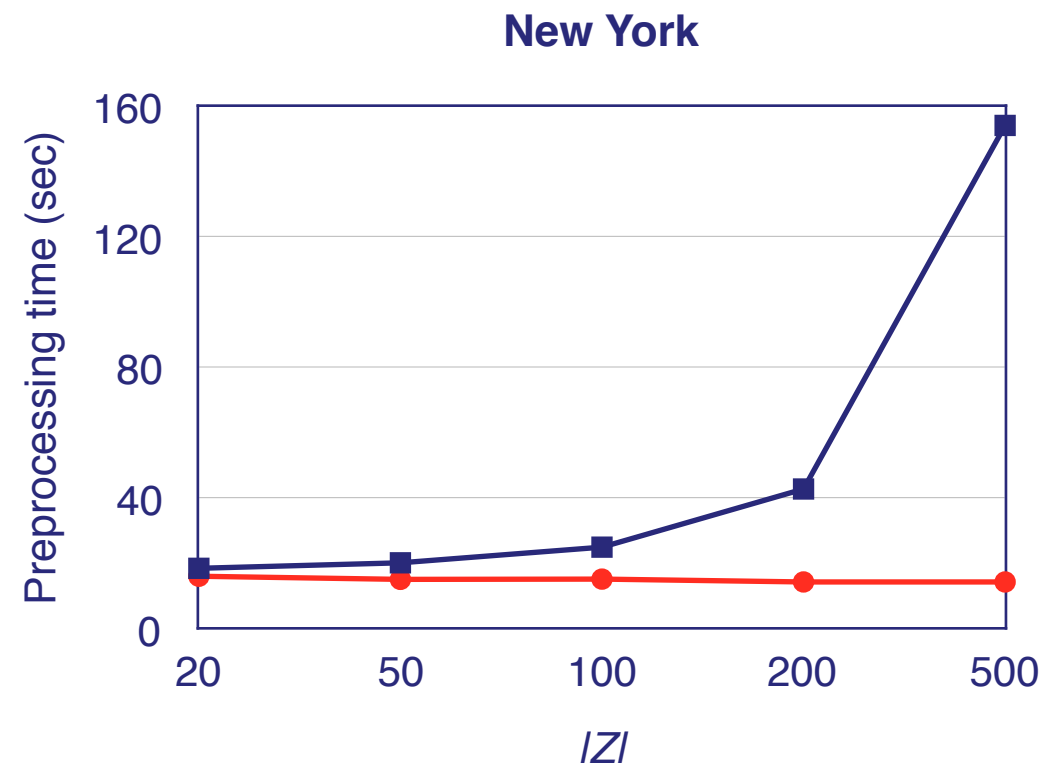
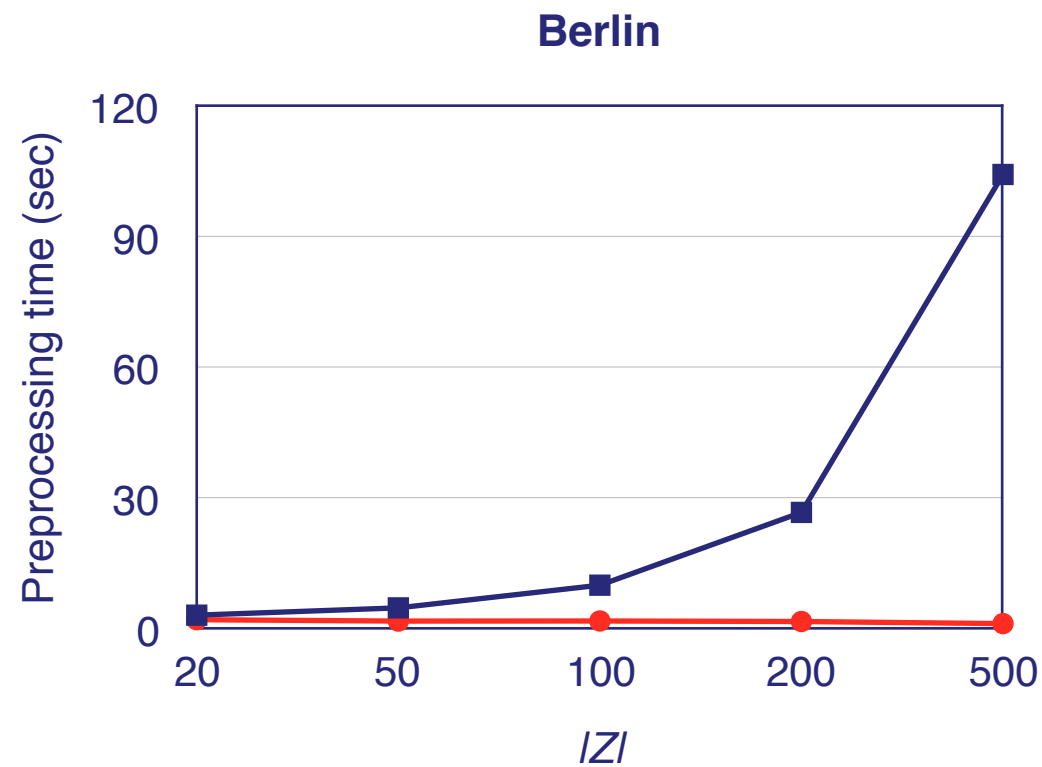


Experimental Evaluation MPUP

- Two datasets:
 - ▶ Berlin (37,126 nodes and 102,260 edges)
 - ▶ New York (264,346 nodes and 730,100 edges)
- Two experiments:
 - ▶ varying number of zones $|Z|$
 - ▶ varying radius of zones r
- Default values: $|Z| = 100$, $r = 1500\text{m}$
- We use Contraction Hierarchies to optimize both preprocessing and query processing

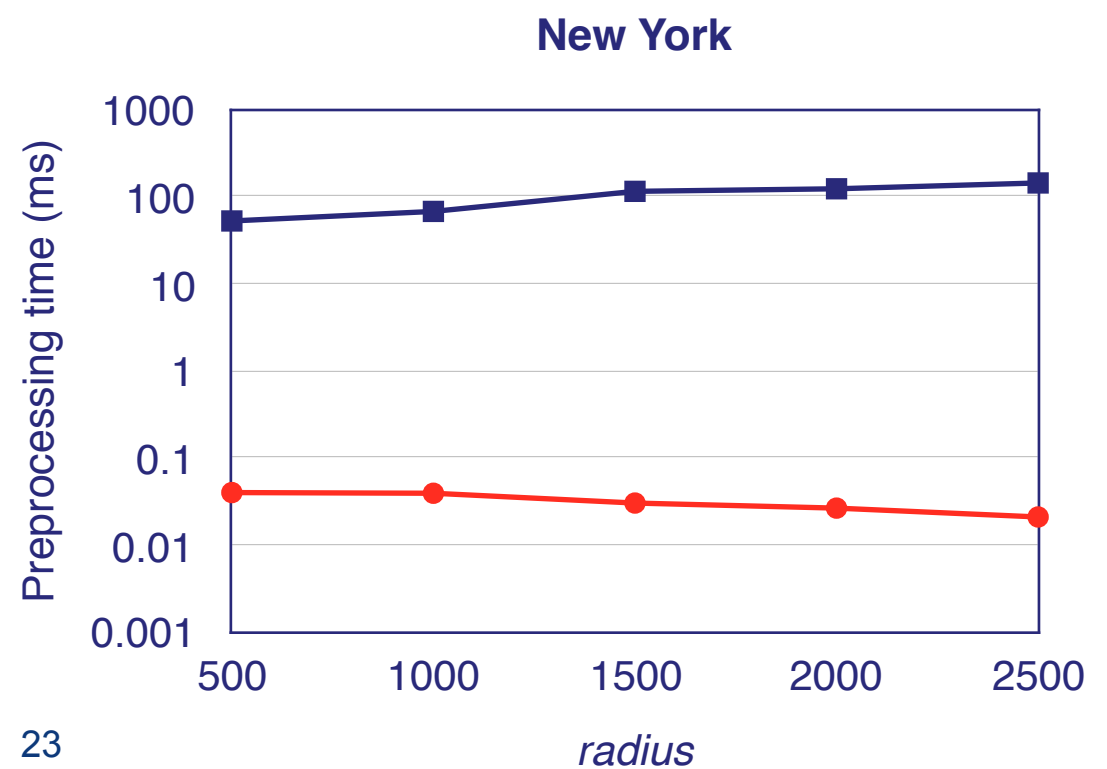
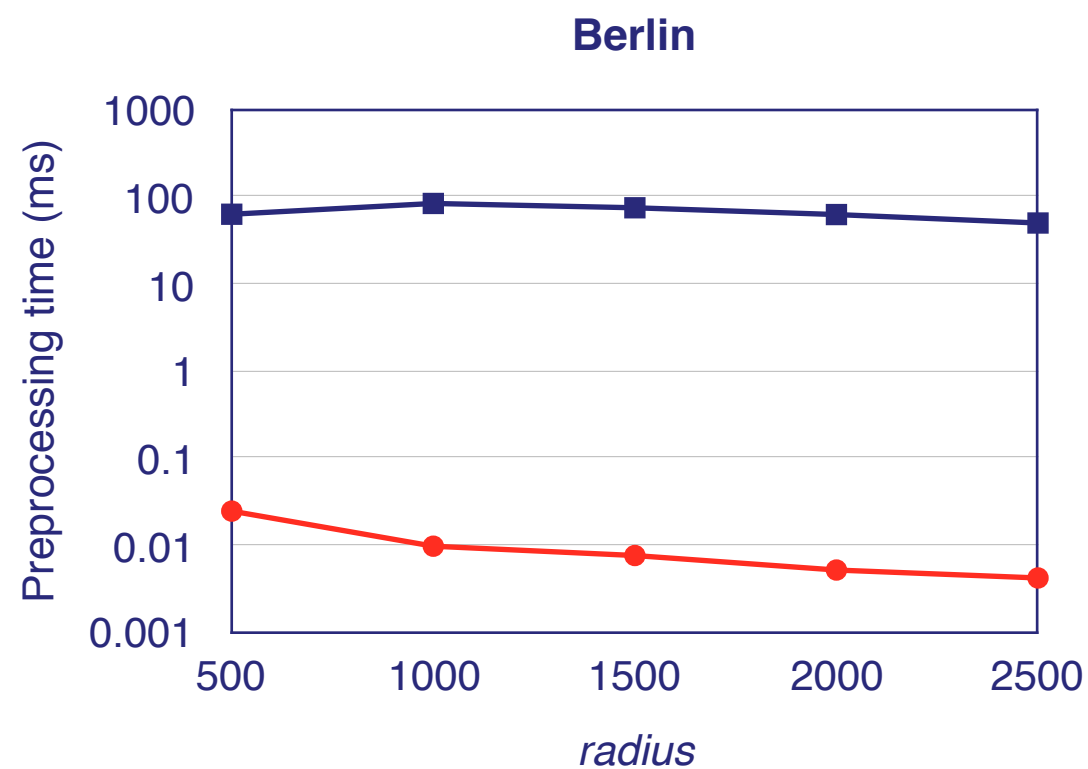
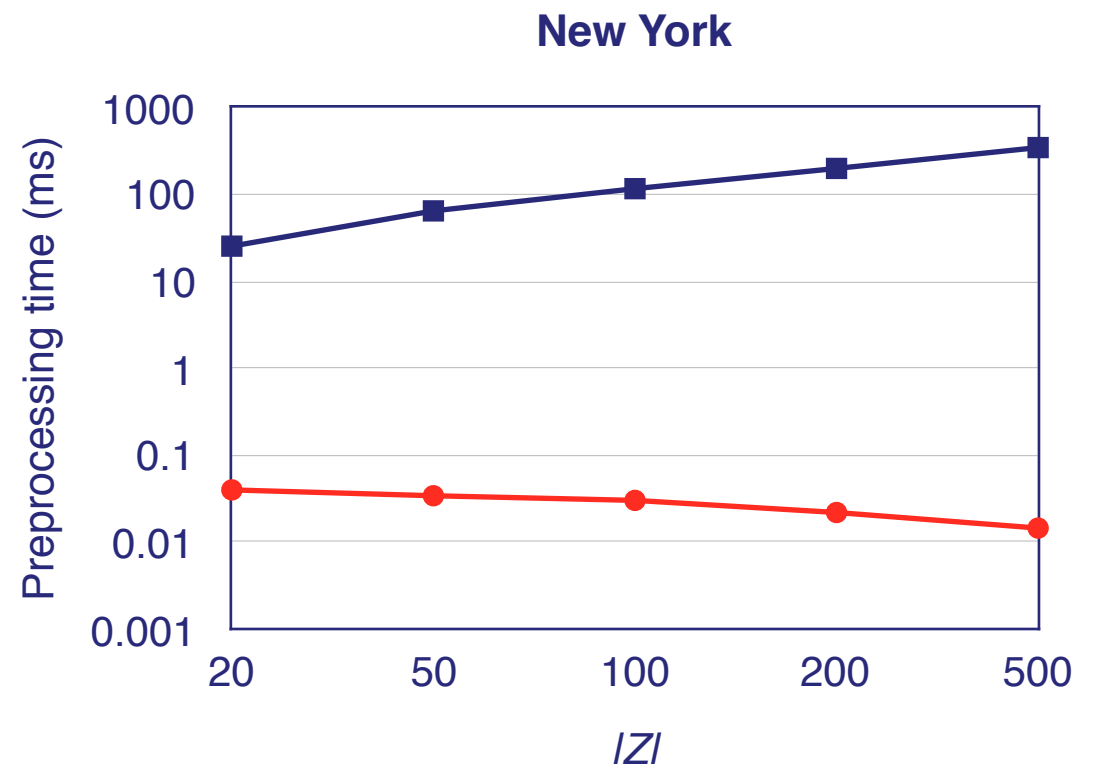
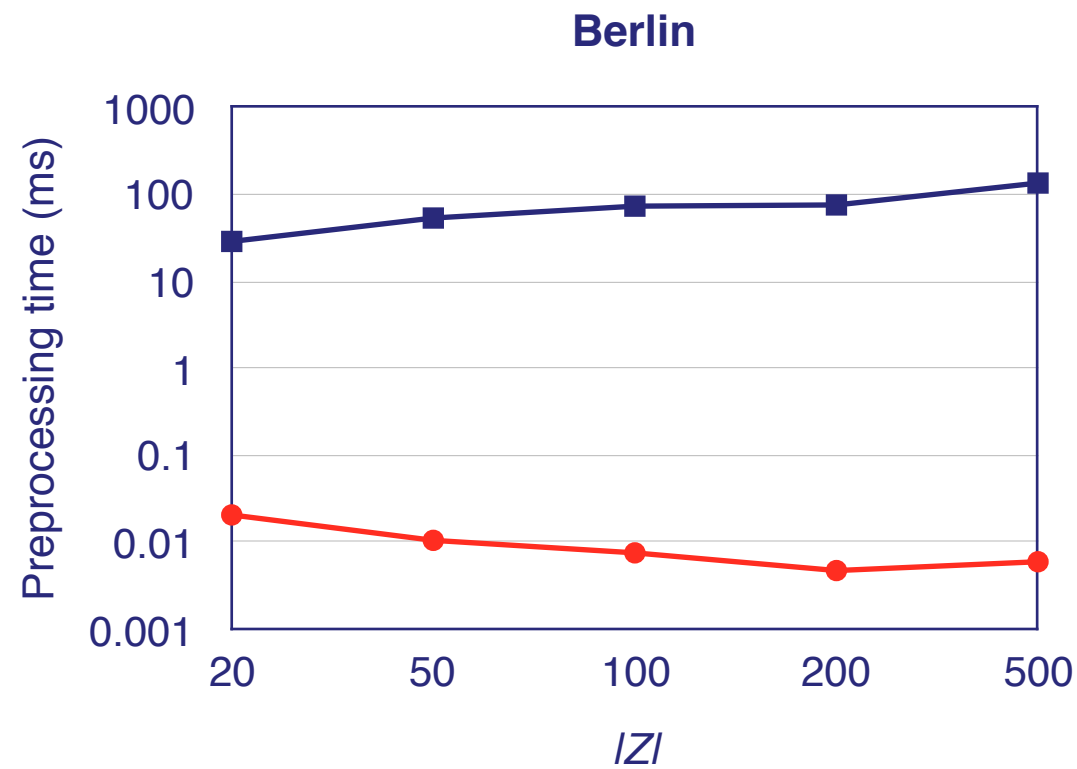
MPUP - Preprocessing

■ Hyperedges ● Compressed Network



MPUP - Query Processing

■ Hyperedges ● Compressed Network



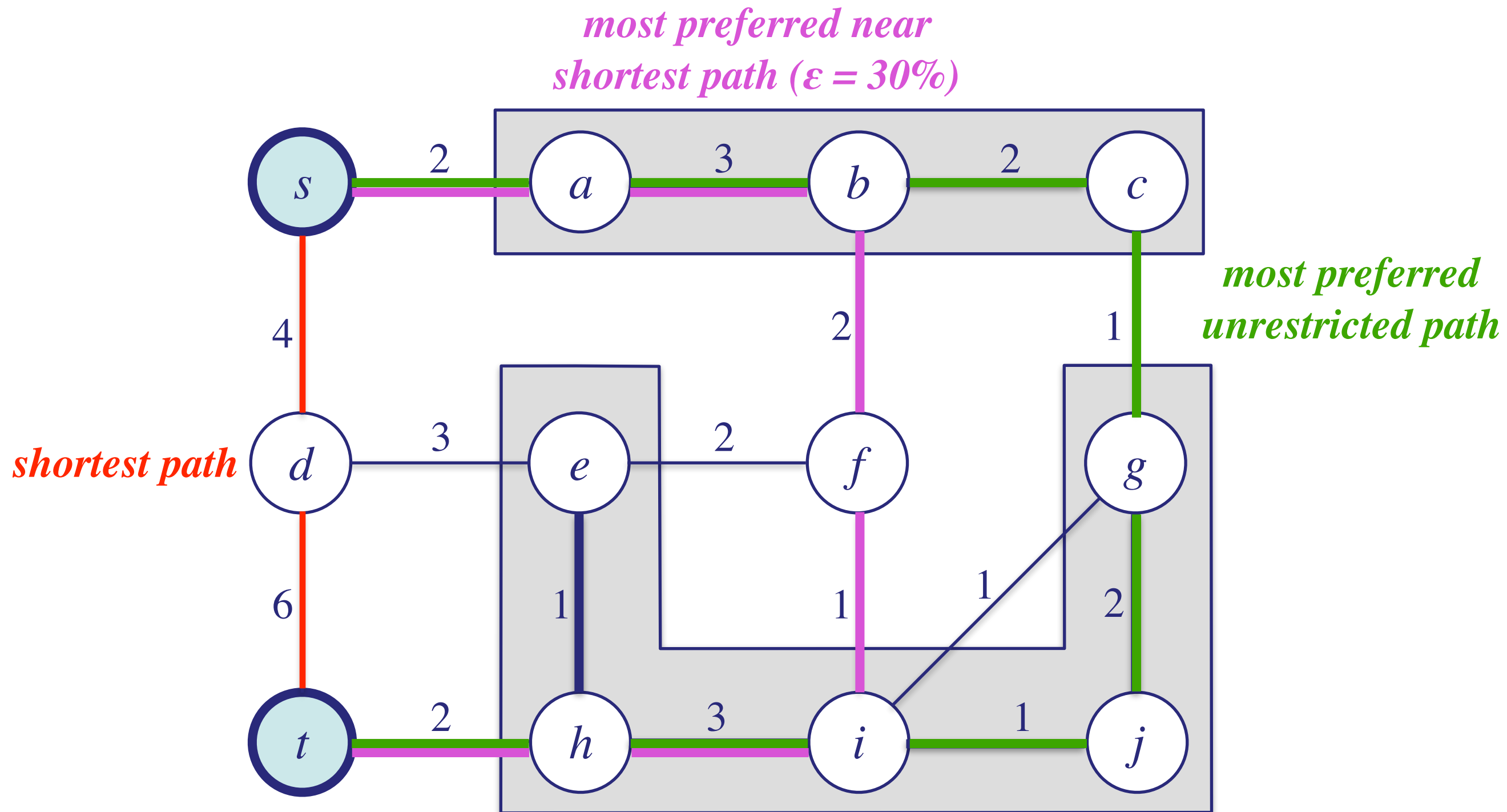
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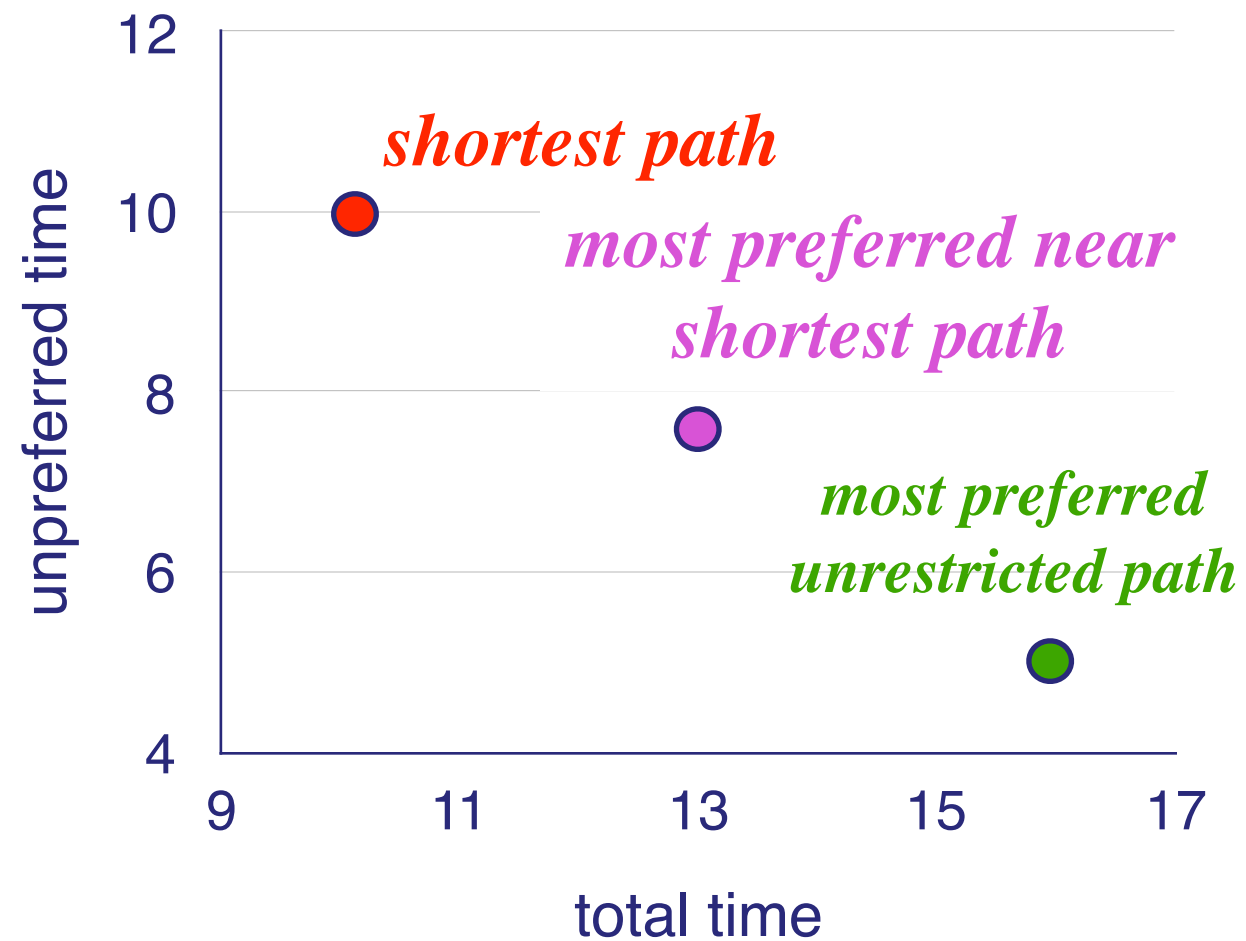
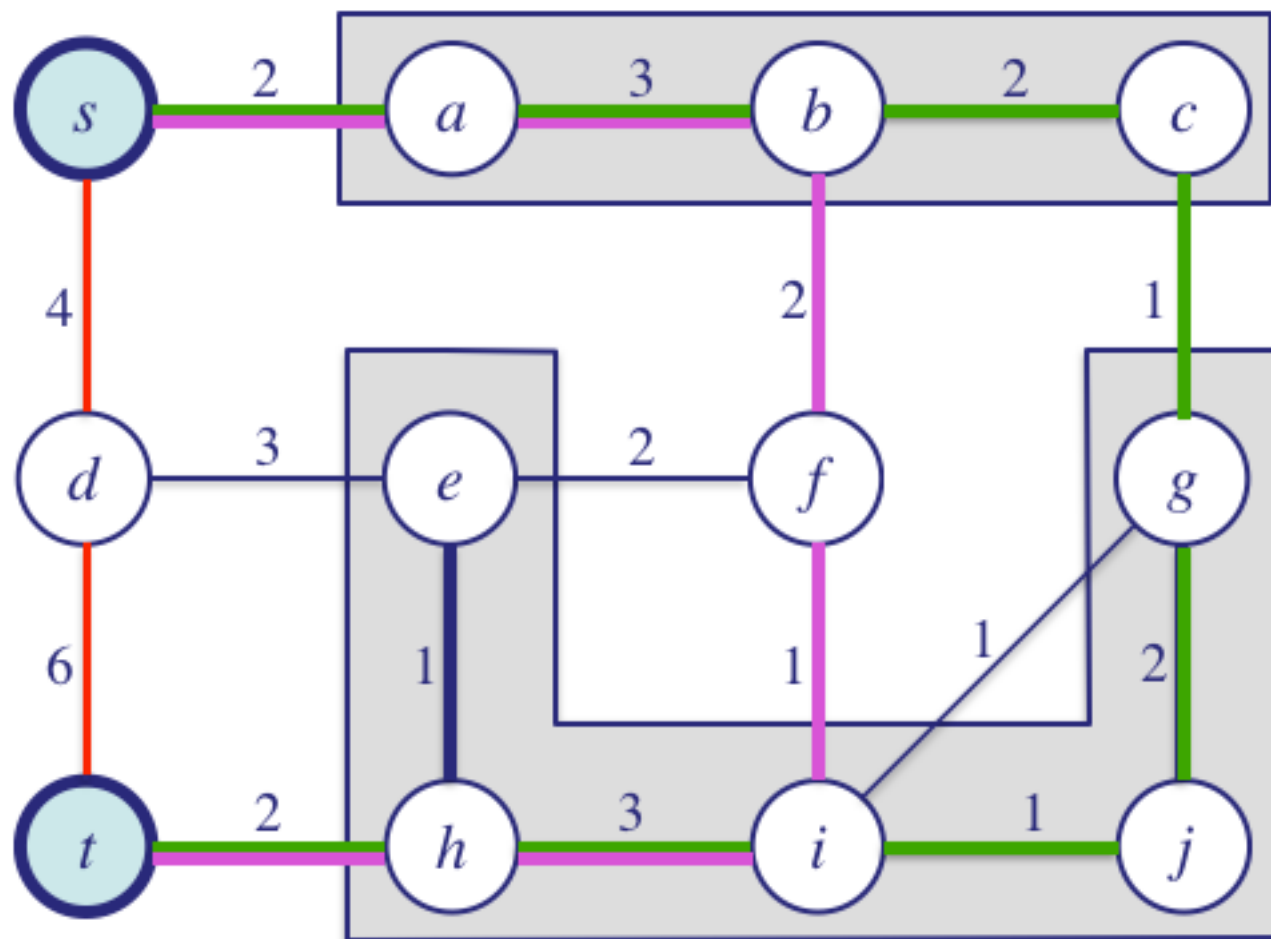
Most Preferred Near Shortest Path

- Find the path which
 - (1) minimizes the time spent outside the Preferred Network
 - (2) is at most $X\%$ longer than the shortest path

Most Preferred Near Shortest Path



Most Preferred Near Shortest Path

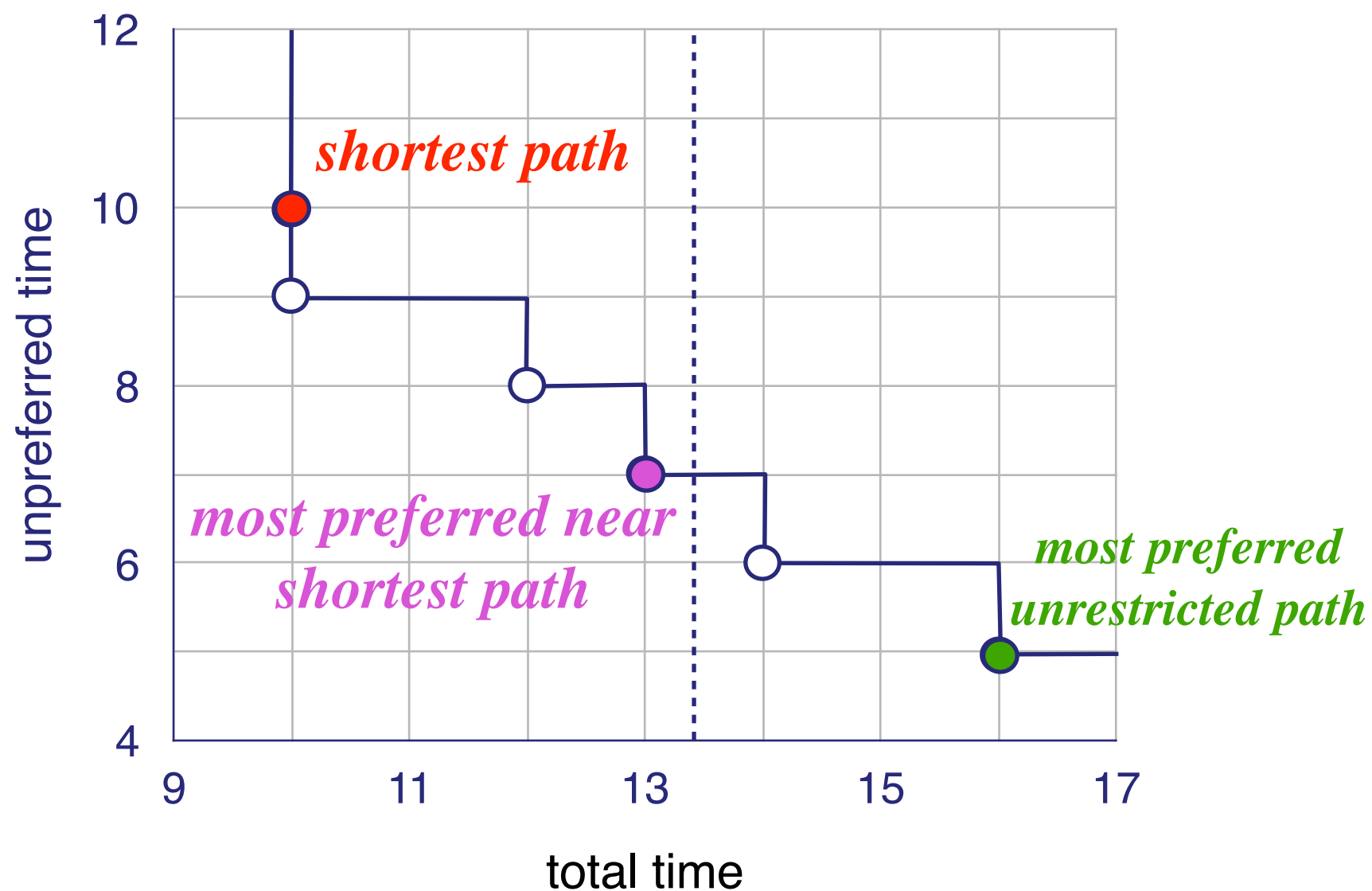


Algorithms

- Advanced Route Skyline Computation (ARSC)

Advanced Route Skyline Computation

- Compute the entire Route Skyline
- Then retrieve the MPNSP

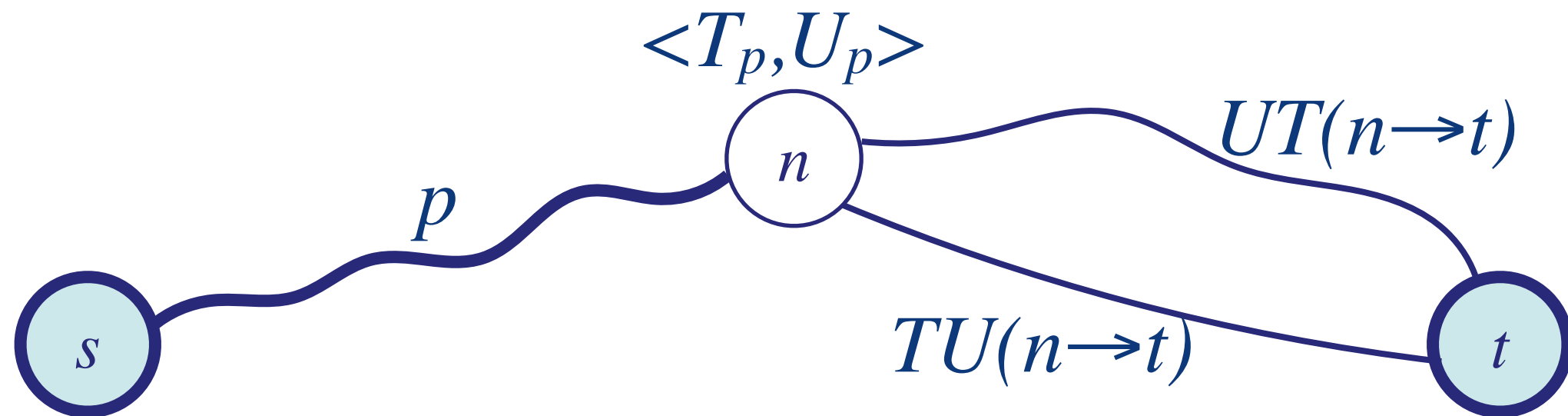


Algorithms

- Advanced Route Skyline Computation (ARSC)
 - ▶ prunes dominated paths
- **Algorithm ALGO-U**
 - ▶ prunes dominated paths
 - ▶ employs upper bounds for the unpreferred time

ALGO-U directs the search towards the desired result of the skyline

Algorithm ALGO-U



Upper Bound U^*

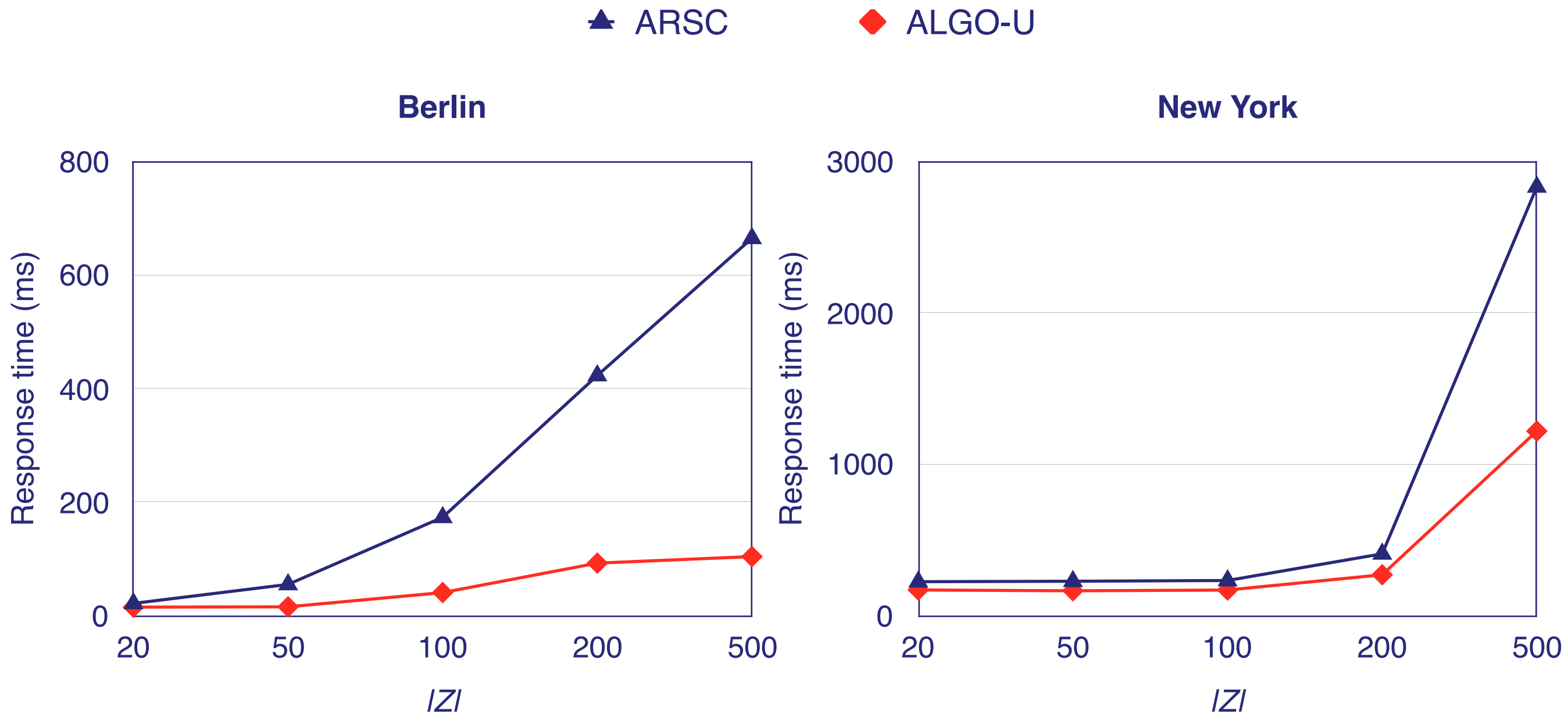
$$\min \begin{cases} U_p + TU(n \rightarrow t) \cdot U, & \text{if } T_p + TU(n \rightarrow t) \cdot T \leq (1 + \varepsilon) \cdot d(s, t) \\ U_p + UT(n \rightarrow t) \cdot U, & \text{if } T_p + UT(n \rightarrow t) \cdot T \leq (1 + \varepsilon) \cdot d(s, t) \end{cases}$$

if the extension does not violate ε

Experimental Evaluation MPNSP

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 - ▶ varying threshold ε
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Experimental Evaluation MPNSP



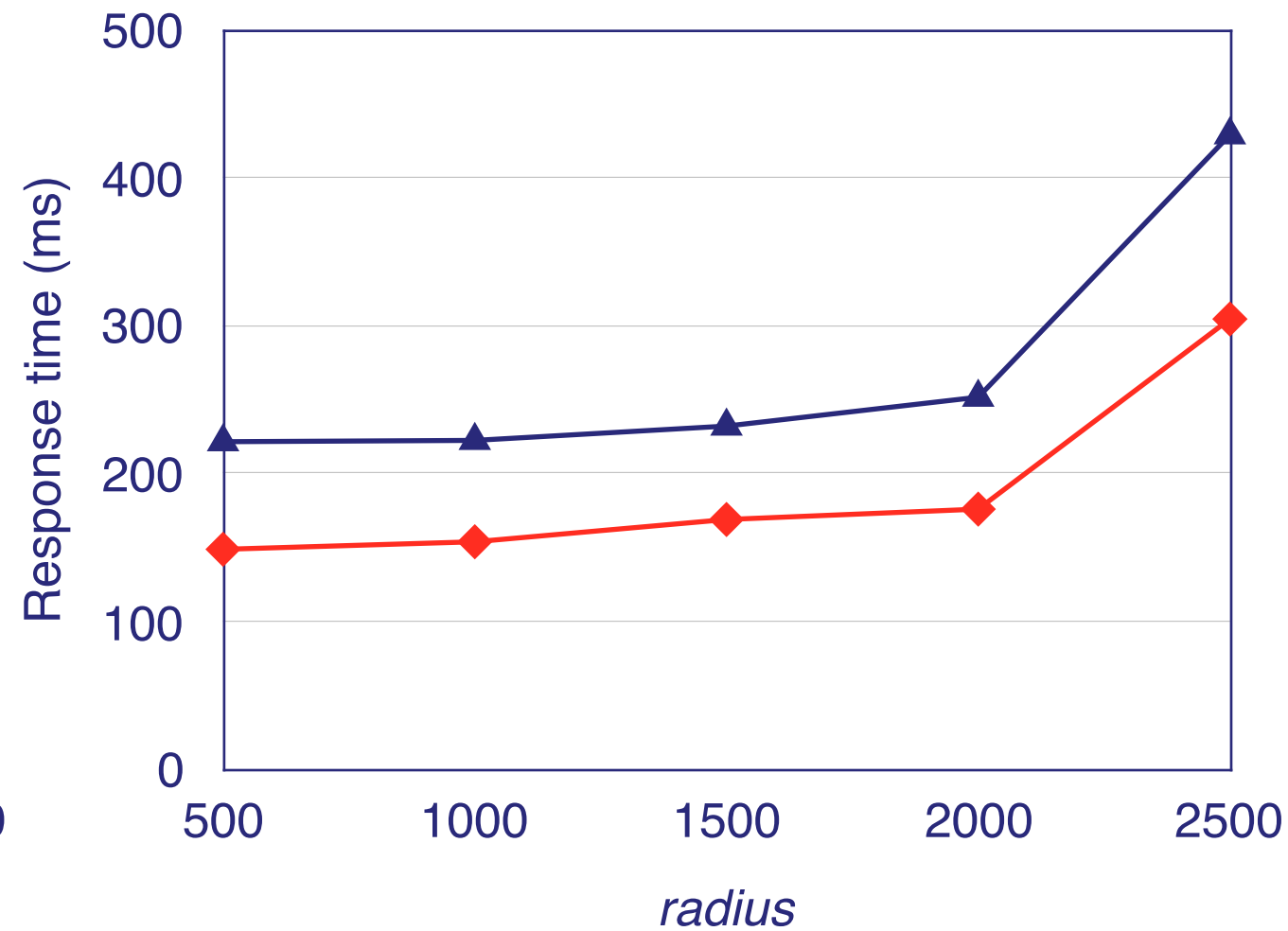
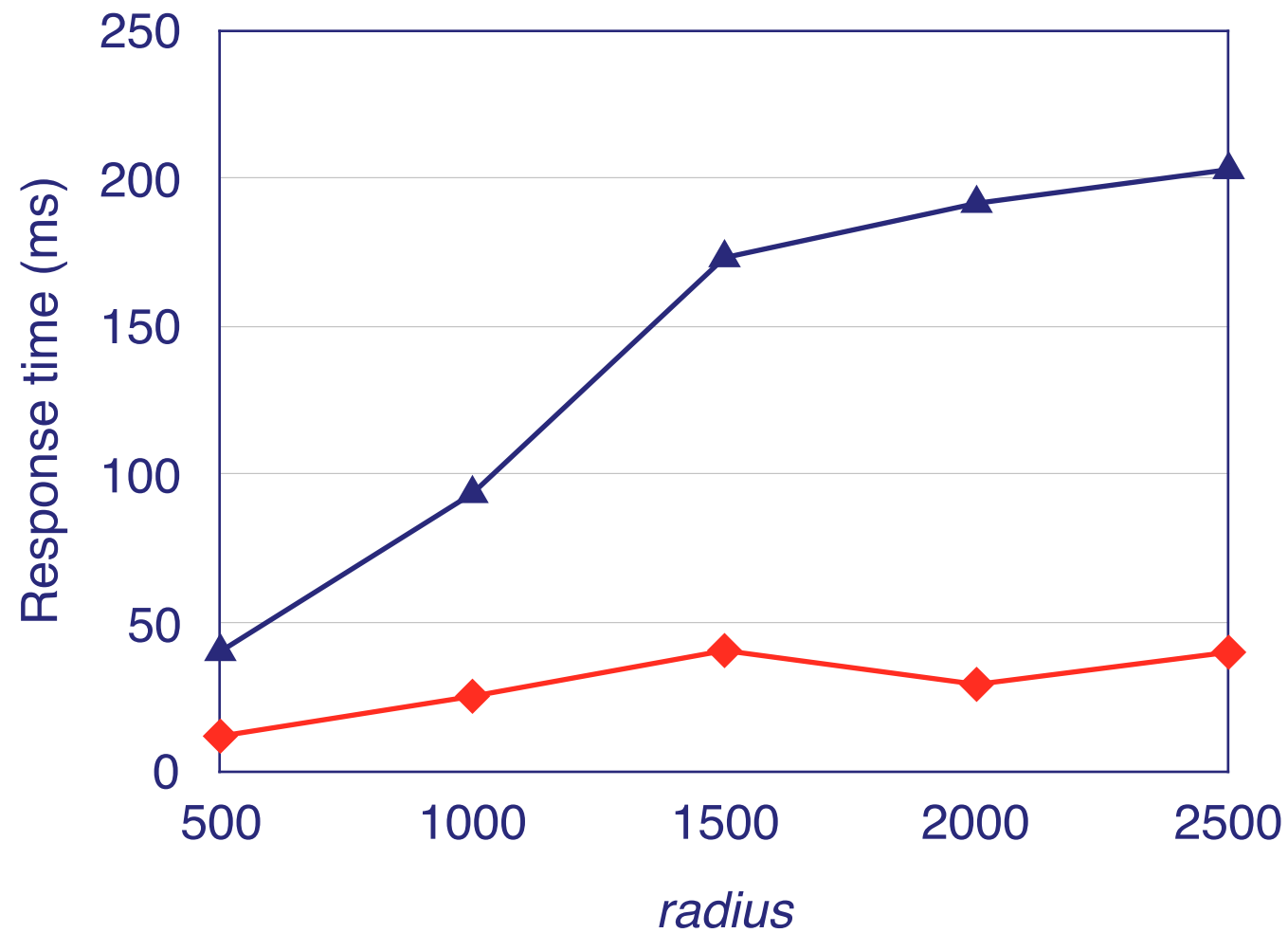
Experimental Evaluation MPNSP

▲ ARSC

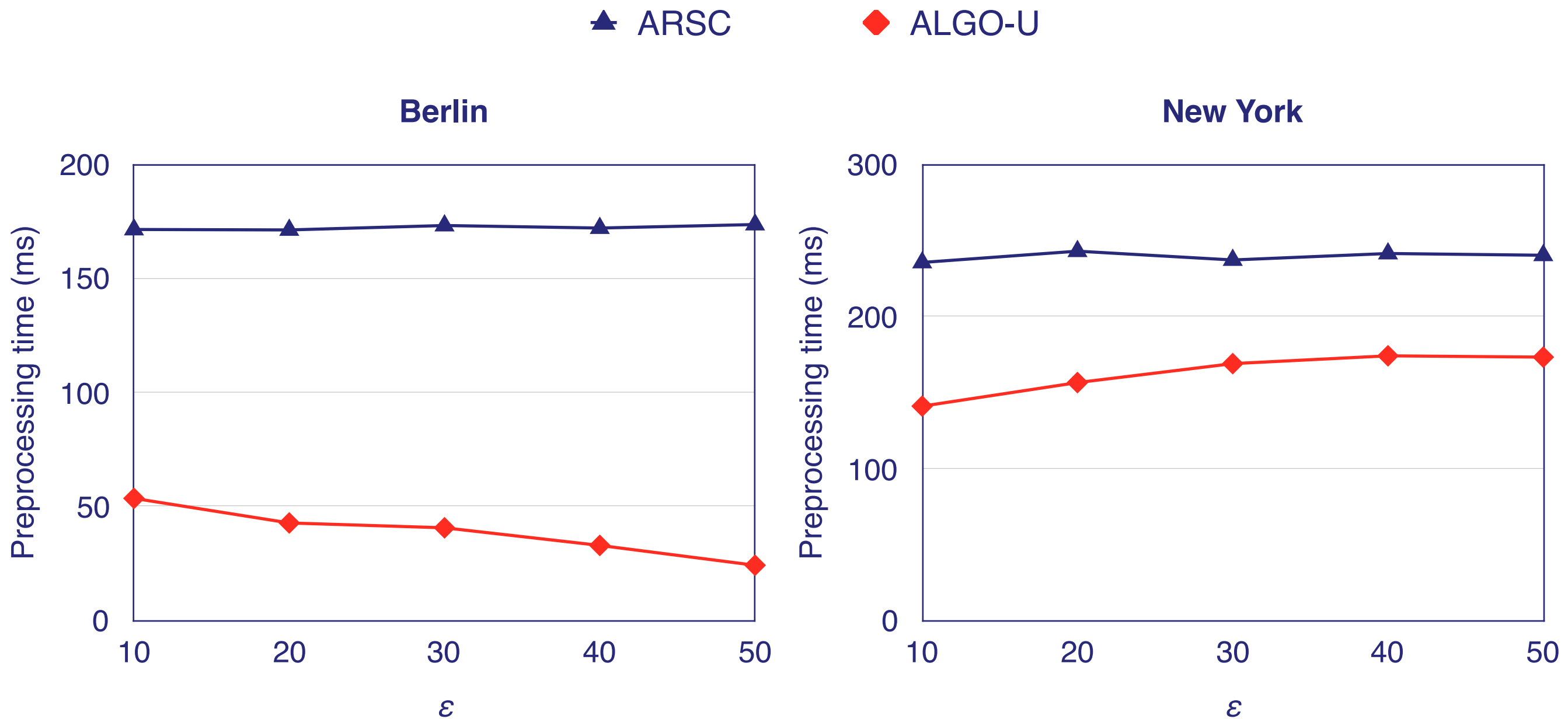
◆ ALGO-U

Berlin

New York



Experimental Evaluation MPNSP



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Conclusions and Future Work

- Conclusions

- ✓ We studied the problem of Finding the Most Preferred Path on road networks
- ✓ The Compressed Network approach improves the state-of-the-art for MPUP
- ✓ We introduced MPNSP along with ALGO-U which also improves the state-of-the-art

- Future Work

- ✓ Investigate pre-processing methods for MPNSP
- ✓ Study methods to extract Preferred Zones



Thank you!